

# ADDENDUM TO “TOPOLOGICAL COMPUTING OF ARRANGEMENTS WITH (CO)CHAINS”

ALBERTO PAOLUZZI

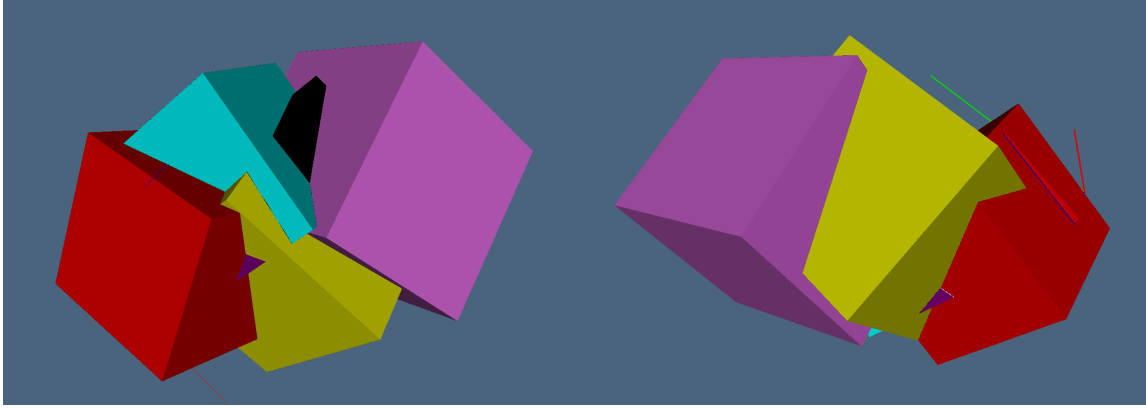


FIGURE 1. Two images of the biggest 3-cell, whose outer 2-boundary coincides with the (reversed) 2-boundary of the outer space (outer 3-cell).

Archimedes to Eratosthenes greeting. ... certain things first became clear to me by a mechanical method, although they had to be demonstrated by geometry afterwards because their investigation by the said method did not furnish an actual demonstration. But it is of course easier, when we have previously acquired by the method, some knowledge of the questions, to supply the proof than it is to find it without any previous knowledge.

*The Method*

## 1. INTRODUCTION

The aims of this short essay are: (a) to give an example of the computational techniques introduced with the paper “Topological computing of arrangements with (co)chains” [1], submitted about two years ago to ACM Transactions on Spatial Algorithms and Systems (TSAS), and (b) to produce some measures providing a benchmark of this approach.

We intend to illustrate here the input and the output data generated by our open-source implementation<sup>1</sup>, when applied to a small collection of closed polyhedral surfaces, so producing the *arrangement* (partition) of 3D space, described by the matrices of linear operators  $\delta_0, \delta_1, \delta_2$  and their transposed operators  $\partial_1, \partial_3, \partial_2$ , which go up and down, respectively, between linear chain spaces  $C_0, C_1, C_2$  generated by the 3D partition.

For problem statement, motivation, definitions, discussions, algorithms, applications, and further examples, the reader is referred to [1]. We discuss in the present paper the Julia's script solving the arrangement problem and the output generated, starting from five randomly oriented and dimensioned cubes around the origin, which are shown in Figure 1. Finally we compare the amount of data produced by this approach with the more crude raw format of graphics applications, i.e. the boundary triangle mesh of generated 3-cells.

## 2. COMPUTATION OF 3D ARRANGEMENT EXAMPLE

**2.1. Partition into solid cells.** A *chain complex* is a short exact sequence of graded linear spaces  $C_p$  of (co)chains, with linear boundary/coboundary maps  $\partial_p$  and  $\delta_p = \partial_{p+1}^\top$ :

$$C_\bullet = (C_p, \partial_p) := C_3 \begin{matrix} \xleftarrow{\delta_2} \\ \xrightarrow{\partial_3} \end{matrix} C_2 \begin{matrix} \xleftarrow{\delta_1} \\ \xrightarrow{\partial_2} \end{matrix} C_1 \begin{matrix} \xleftarrow{\delta_0} \\ \xrightarrow{\partial_1} \end{matrix} C_0.$$

The cells of a space partition are one-to-one with the basis elements of chain spaces. The chain maps ( $\partial_p$  or  $\delta_p$ ) fully describe the topology of the cellular arrangement.

The set of manifold 3-cells of the 3D arrangement produced by the surfaces in Figure 1 is shown in Figure 2. It is worthwhile to note that such solid cells may be non-convex and non-contractible, i.e., non simply connected. Such properties extend to their 2-cell faces.

It may be interesting to notice also that cells may have different topological genres, i.e., any number of holes and tunnels. Look for this purpose at the fourth cell (from left) of the second row (from top) of Figure 2. Other 3-cells have also smaller holes in the faces and/or tunnels within. The boundary triangulation of 3-cells needed to generate graphics on a display device or for 3D printing is also provided in our package repository on [github.com](https://github.com).

The geometric model of the arrangement is generated as a pair (*Geometry*, *Topology*), where the first one is simply given by the embedding map  $V$  of vertices (0-cells), i.e. by their coordinate matrix, given by columns:

```
julia> V
3x137 Array{Float64,2}:
 1.01181  0.215639  0.516927  0.449016 ... 1.01221   1.30039   0.741732  1.02991
 0.160033 0.06801   0.102833  0.094984 ... 0.816151  0.545669  1.51776   1.24728
 0.196256 0.206963  0.202911  0.203825 ... 0.249344  0.985248  0.613139  1.34904
```

The cardinality of 0-, 1-, 2-, and 3-cell sets  $V$ ,  $E$ ,  $F$ ,  $C$  are given as row and column numbers of the sparse matrices of their signed incidence relations  $EV$ ,  $FE$ ,  $CF$ , that can be interpreted as coboundary operators  $\delta_0$  (from vertices to edges),  $\delta_1$  (from edges to faces), and  $\delta_2$  (from faces to solid cells), respectively. The prefix `cop` stands for *chain operator*:

<sup>1</sup>See our [Github.com](https://github.com) repository,

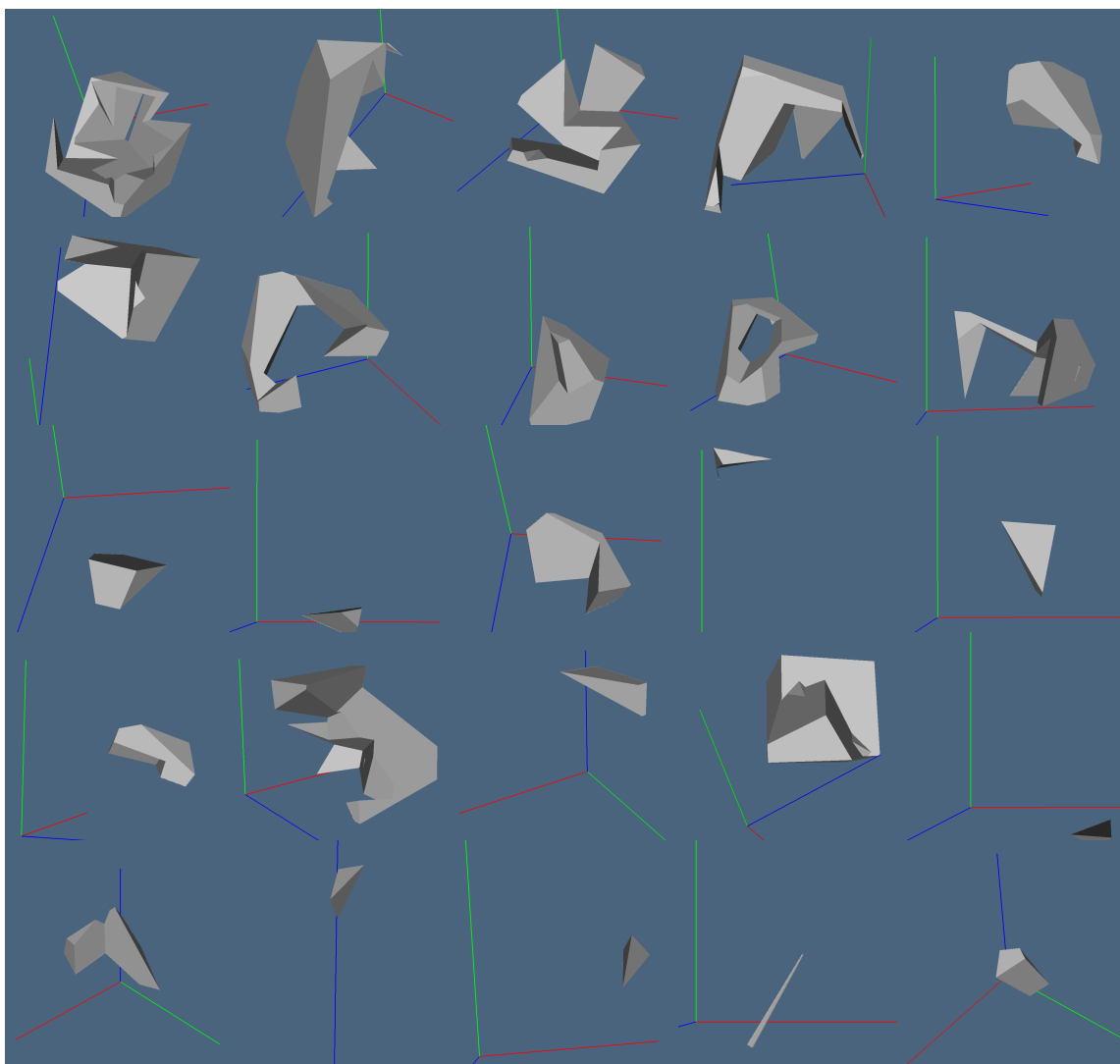


FIGURE 2. The 3-cells of the arrangement of  $\mathbb{E}^3$  generated by a collection of five random 3-cubes. Their assembly provides the union of the five 3-cubes. Each 3-cell is given by a column of the sparse matrix of map  $\delta_2 : C_2 \rightarrow C_3$ .

julia> copEV	julia> copFE	julia> copCF
268x137 SparseMatrixCSC	157x268 SparseMatrixCSC	27x157 SparseMatrixCSC
with 536 stored entries:	with 786 stored entries:	with 314 stored entries:
[1 , 1] = -1	[1 , 1] = 1	[1 , 1] = -1
[5 , 1] = -1	[11 , 1] = 1	[19 , 1] = 1
[39 , 1] = -1	[2 , 2] = 1	[2 , 2] = -1
...	...	...
[258, 137] = 1	[156, 267] = 1	[21 , 156] = 1
[260, 137] = 1	[150, 268] = 1	[1 , 157] = 1
[268, 137] = 1	[157, 268] = 1	[21 , 157] = -1

**2.2. The input collection of surfaces.** The input data for the example of Figure 1 was generated by a little script in Julia language, which is not of interest here. Conversely, we provide below the full input, i.e. both the geometry  $V$  and the topology  $EV$ ,  $FV$ . Such text data amount to 2.1k bytes.

```
julia> @show V;
V = [1.01181 0.215639 0.91979 0.123616 1.02252 0.226347 0.930498 0.134324 0.0458309
-0.301827 0.348275 0.0006172 0.579367 0.23171 0.881811 0.534154 -0.0521776 0.627953
-0.190635 0.489496 -0.0233884 0.656742 -0.161846 0.518285 0.27652 -0.0875132 0.52527
0.161237 0.509324 0.145291 0.758074 0.394041 0.27631 0.564484 0.0058279 0.294002
1.01221 1.30039 0.741732 1.02991; 0.160033 0.0680099 0.956278 0.864255 0.160649
0.0686266 0.956895 0.864872 -0.200245 0.102199 0.417839 0.720283 -0.35354 -0.0510965
0.264543 0.566987 0.682359 0.543901 0.0592036 -0.0792537 0.956374 0.817917 0.333219
0.194762 -0.102028 0.146722 0.324834 0.573584 -0.16916 0.0795901 0.257702 0.506452
0.452356 0.181874 1.15396 0.883481 0.816151 0.545669 1.51776 1.24728; 0.196256
0.206963 0.196872 0.20758 0.997729 1.00844 0.998346 1.00905 0.0677451 0.601282
-0.0855504 0.447986 0.502301 1.03584 0.349005 0.882542 0.159301 0.18809 0.433316
0.462105 0.797002 0.825792 1.07102 1.09981 0.1446 0.377404 0.0774682 0.310272 0.580364
0.813168 0.513232 0.746036 0.403805 1.13971 0.767599 1.5035 0.249344 0.985248 0.613139
1.34904]

julia> @show FV,EV;
(FV,EV) = (Array{Int64,1}[[1,2,3,4],[5,6,7,8],[1,2,5,6],[3,4,7,8],[1,3,5,7],[2,4,6,8],
[9,10,11,12],[13,14,15,16],[9,10,13,14],[11,12,15,16],[9,11,13,15],[10,12,14,16],[17,
18,19,20],[21,22,23,24],[17,18,21,22],[19,20,23,24],[17,19,21,23],[18,20,22,24],[25,26,
27,28],[29,30,31,32],[25,26,29,30],[27,28,31,32],[25,27,29,31],[26,28,30,32],[33,34,35,
36],[37,38,39,40],[33,34,37,38],[35,36,39,40],[33,35,37,39],[34,36,38,40]],
Array{Int64,1}[[1,2],[3,4],[5,6],[7,8],[1,3],[2,4],[5,7],[6,8],[1,5],[2,6],[3,7],[4,8],
[9,10],[11,12],[13,14],[15,16],[9,11],[10,12],[13,15],[14,16],[9,13],[10,14],[11,15],
[12,16],[17,18],[19,20],[21,22],[23,24],[17,19],[18,20],[21,23],[22,24],[17,21],[18,
22],[19,23],[20,24],[25,26],[27,28],[29,30],[31,32],[25,27],[26,28],[29,31],[30,32],
[25,29],[26,30],[27,31],[28,32],[33,34],[35,36],[37,38],[39,40],[33,35],[34,36],[37,
39],[38,40],[33,37],[34,38],[35,39],[36,40]])
```

**2.3. The generating script.** Below we compute the cellular 3-complex of the  $\mathbb{E}^3$  space partition induced by the above collection of surfaces. In the current prototype implementation, some transformation of input data format is needed. A simpler API will be provided soon. In particular, type `Lar.Cells` (array of arrays of integers) is converted to type `Lar.ChainOp`, i.e., the Julia's type `SparseMatrixCSC{Int8,Int64}` for sparse matrices.

```
cop_EV = Lar.coboundary_0(EV::Lar.Cells);
cop_EW = convert(Lar.ChainOp, cop_EV);
cop_FE = Lar.coboundary_1(V, FV::Lar.Cells, EV::Lar.Cells);
W = convert(Lar.Points, V');

V, copEV, copFE, copCF = Lar.Arrangement.spatial_arrangement(
    W::Lar.Points, cop_EW::Lar.ChainOp, cop_FE::Lar.ChainOp )
```

**2.4. The output chain complex.** It is easier to compare different data structures by their data object numbers and file representation, without taking into account the actual computational architecture and the memory occupancy.

Therefore, some comparison numbers follow: into the five (unconnected) boundaries of the input `Lar` cubes there are  $5 \times 8 = 40$  vertices,  $5 \times 6 = 30$  quads, and  $5 \times 12 = 60$  edges. For the computed space arrangement, 137 vertices and 953 triangles are contained within the 27 cell groups of the output `OBJ` file. The arrays of vertices, edges, faces and 3-cells, without any added data structure<sup>2</sup>, are given in Ref.<sup>3</sup> and weight for 38 kB as text file.

The cellular 3-complex output with 137 0-cells, 268 1-cells, 157 2-cells, 27 3-cells, for a total of 589 graded  $p$ -cells, is better represented as a *geometric chain complex*. In particular, the textual values of `V`, of sparse matrices `copEV`, and `copFE`, computed by the *Merge* algorithm [1], plus the sparse matrix `copCF` computed by the *TGW* algorithm [1] in 3D, are given in `github.com`/<sup>4</sup> for a file size of  $(16+2)k$  bytes.

It might be interesting to note that the minimal and fairly crude `OBJ` file representation of 3-cells, also provided in <sup>5</sup>, amounts to 15k bytes. But whereas this last format just contains a bunch of triangles as triples of indices of vertices, without any storage of the assembly and cells topology, the *chain complex* representation conversely allows for a complete understanding of topology and the direct querying of either single or multiple topological features, using only sparse matrix-vector or sparse matrix-matrix multiplications.

### 3. CONCLUSION

In this short *Addendum to Topological Computing of Arrangements with (Co)chains*, a small example of computation of a 3D arrangement has been presented and discussed. Our main aim was to show the fairly general nature of generated 3-cells, that are connected and non manifold, non contractible, and with any topological genus. We would like to remark also the briefness and simplicity of this representation of *geometric models* as *chain complexes*, and its great generality. This author hopes that the presented material may help to better understand such a novel approach to geometric computing using algebraic topological tools.

### REFERENCES

1. Alberto Paoluzzi, Vadim Shapiro, Antonio DiCarlo, Francesco Furiani, Giulio Martella, and Giorgio Scorzelli, *Topological computing of arrangements with (co)chains*, submitted to Transactions on Spatial Algorithms and Systems, ACM, New York, NY (August 2017).

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<sup>2</sup>The incidence relations are computable on demand by product or transposition of sparse binary *characteristic matrices*, as discussed in Reference [1]

<sup>3</sup><https://github.com/cvdlab/LinearAlgebraicRepresentation.jl/tree/julia-1.0/examples/3d/data>,

<sup>4</sup>*idem*

<sup>5</sup>*idem*